

## Valid Well-Formed Formulae of Ordinary Propositional Calculus

**FR1** A letter standing alone is a wff

**FR2** If  $\alpha$  is a wff, so is  $\sim \alpha$

**FR3** If  $\alpha$  and  $\beta$  are wff, so is  $\alpha \vee \beta$

**Def**  $\wedge (\alpha \wedge \beta) =_{df} \sim (\sim \alpha \vee \sim \beta)$

**Def**  $\supset (\alpha \supset \beta) =_{df} (\sim \alpha \vee \beta)$

**Def**  $\equiv (\alpha \equiv \beta) =_{df} (\alpha \supset \beta) \wedge (\beta \supset \alpha)$

**PC1**  $(p \wedge q) \supset p$

**PC2**  $(p \wedge q) \supset q$

**PC3**  $(p \supset q) \supset ((p \supset r) \supset (p \supset (q \wedge r)))$  (Comp)

**PC4**  $p \supset (q \supset (p \wedge q))$  (Adj)

**PC5**  $(p \supset q) \supset ((q \supset p) \supset (p \equiv q))$

**PC6**  $(p \supset q) \supset ((q \supset r) \supset (p \supset r))$  (Syll)

**PC7**  $(p \supset (q \supset r)) \supset ((p \wedge q) \supset r)$  (Imp)

**PC8**  $(p \supset q) \supset ((q \supset (r \supset s)) \supset ((p \wedge r) \supset s))$

**PC9**  $p \supset (p \vee q)$

**PC10**  $q \supset (p \vee q)$

**PC11**  $(p \supset q) \supset ((r \supset q) \supset ((p \vee r) \supset q))$

**PC12**  $p \equiv \sim \sim p$  (DN)

**PC13**  $(p \vee q) \equiv \sim (\sim p \wedge \sim q)$  (DeM)

**PC14**  $(p \wedge q) \equiv \sim (\sim p \vee \sim q)$  (DeM)

**PC15**  $(p \supset q) \equiv (\sim q \supset \sim p)$  (Transp)

**PC16**  $(p \vee q) \equiv (q \vee p)$  (Comm)

**PC17**  $(p \wedge q) \equiv (q \wedge p)$  (Comm)

**PC18**  $((p \vee q) \vee r) \equiv (p \vee (q \vee r))$  (Assoc)

**PC19**  $((p \wedge q) \wedge r) \equiv (p \wedge (q \wedge r))$  (Assoc)

**PC20**  $p \equiv (p \vee p)$

**PC21**  $p \equiv (p \wedge p)$

## Modal Logic

In addition to the rules and formulae of PC:

**FR4** If  $\alpha$  is a wff, so is  $L\alpha$

**Def M**  $M\alpha =_{df} \sim L \sim \alpha$

**K**  $L(p \supset q) \supset (Lp \supset Lq)$  (Necessarily)

**K1**  $L(p \wedge q) \supset (Lp \wedge Lq)$

**K2**  $(Lp \wedge Lq) \supset L(p \wedge q)$

**K3**  $L(p \wedge q) \equiv (Lp \wedge Lq)$

**K4**  $Lp \vee Lq \supset L(p \vee q)$

**K5**  $Lp \equiv \sim M \sim p$

**K6**  $M(p \vee q) \equiv (Mp \vee Mq)$

**K7**  $M(p \supset q) \equiv (Lp \supset Mq)$

**K8**  $M(p \wedge q) \supset (Mp \wedge Mq)$

**K9**  $L(p \vee q) \supset (Lp \vee Mq)$

**DR1**  $\vdash \alpha \supset \beta \rightarrow \vdash L\alpha \supset L\beta$

**DR2**  $\vdash \alpha \equiv \beta \rightarrow \vdash L\alpha \equiv L\beta$

**DR3**  $\vdash \alpha \supset \beta \rightarrow \vdash M\alpha \supset M\beta$

**D** is the extension of **K** in which every world sees at least one world (serial)

**D**  $Lp \supset Mp$

**D1**  $M(p \supset p)$

**P**  $\vdash M\alpha \rightarrow \vdash \alpha$

**T** is the extension of **K** in which every world sees itself (reflexive)

**T**  $Lp \supset p$

**T1**  $p \supset Mp$

**T2**  $M(p \supset Lp)$

**S4** is the extension of **T** (reflexive) with transitivity; **S4** has 14 distinct modalities

**S4**  $Lp \supset LLp$

**S4-1**  $MMp \supset Mp$

**S4-2**  $Lp \equiv LLp$

**S4-3**  $Mp \equiv MMp$

**S4-4**  $MLMp \supset Mp$

**S4-5**  $LMp \supset LMp$

**S4-6**  $LMp \equiv LMp$

**S4-7**  $MLP \equiv MLMLp$

**B** (Borruwerian system) is an extension of **T** (reflexive) which is symmetric.

**B**  $p \supset LMp$

**B1**  $MLp \supset p$

**DR4**  $\vdash M\alpha \supset \beta \rightarrow \vdash \alpha \supset L\beta$

**S5** is the extension of **T** (reflexive) with transivity (**S4**) and symmetry (**B**); **S5** has 6 distinct modalities

**S5**  $Mp \supset LMp$

**S5-1**  $MLp \supset Lp$

**S5-2**  $Mp \equiv LMp$

**S5-3**  $Lp \equiv MLp$

**S4**  $Lp \supset LLp$

**S5-4**  $L(p \vee Lq) \equiv (Lp \vee Lq)$

**S5-5**  $L(p \vee Mq) \equiv (Lp \vee Mq)$

**S5-6**  $M(p \wedge Mq) \equiv (Mp \wedge Mq)$

**S5-7**  $M(p \wedge Lq) \equiv (Mp \wedge Lq)$